

Approximate Stream Reasoning with Incomplete State Information

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Introduction

- 1 Introduction
- 2 Stream Reasoning with Incomplete Information
- 3 Progression Graph-Based Progression
- 4 Summary

Introduction

Consider runtime verification for checking whether an agent is behaving in a safe manner.

Example (Safety)

“A robot in an unsafe state should return to a safe state within 10 seconds”



Motivation: Incomplete information!

Metric Temporal Logic

We use **Metric Temporal Logic** (MTL) as a language for describing temporal rules that must hold.

Definition (MTL syntax)

The syntax for MTL is as follows for atomic propositions $p \in \text{Prop}$, temporal interval $I \subseteq (0, \infty)$, and well-formed formulas (wffs) ϕ and ψ :

$$p \mid \neg\phi \mid \phi \vee \psi \mid \phi \mathcal{U}_I \psi$$

where \Box_I and \Diamond_I are syntactic sugar for ‘always’ and ‘eventually’.

Progression-based Runtime Verification

Progression is an incremental **syntactic rewriting procedure** introduced by Bacchus and Kabanza (1996, 1998):

MTL Formula + Complete State + Delay \Rightarrow MTL Formula

$$\frac{\phi_0 = \Box(\neg p \rightarrow \Diamond_{[0,10]} p), s = \{\neg p\}, \Delta = 2}{\phi_1 = \Diamond_{[0,8]} p \wedge \Box(\neg p \rightarrow \Diamond_{[0,10]} p)}$$

Stream Reasoning with Incomplete Information

Problem: How to perform progression with **incomplete** states?

General idea: Apply model counting

Incomplete States and Streams

Important assumptions:

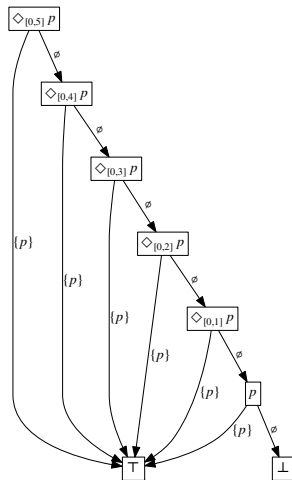
- We keep a constant delay value (Δ) and omit it from here on;
- An **incomplete state** \hat{s} is a disjunctive set of complete states;
- A (piecewise) **incomplete stream** $\hat{\rho}$ is a sequence of incomplete states;
- We assume we have a probabilistic model of a stream denoted by a **state universe** \mathbf{S}_n for every time-point n .

Progression Graphs

A **progression graph** encodes formulas and their progressions into a graph

$G(\chi) = (\chi, V, E)$ such that

- vertices represent formulas;
- $\chi \in V$ represents the graph source formula; and
- labelled edges $(\phi, \psi, s) \in E$ iff $\text{PROGRESS}(\phi, s) = \psi$.



Progression Graphs

Progression graphs $G_n(\chi) = (\chi, V, E, m_n)$ carry **probability mass**:

$$m_0(\chi) = 1.0 \text{ (Initialization)}$$

$$m_n(v) = \sum_{(v', v, s) \in E} (m_{n-1}(v') \Pr[\mathbf{S}_n = s \mid \hat{\mathbf{S}}_n])$$

Theorem (Soundness)

Given a progression graph $G_n(\chi)$ and a stream $\hat{\rho}$:

$$\lim_{n \rightarrow \infty} m_n(\top) = \Pr[\hat{\rho}, t_0 \models \chi],$$

$$\lim_{n \rightarrow \infty} m_n(\perp) = \Pr[\hat{\rho}, t_0 \not\models \chi].$$

Progression Graph-Based Progression

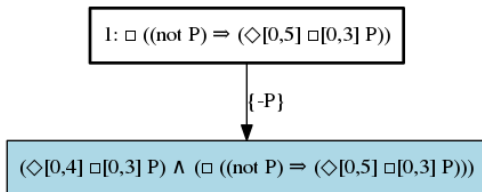
Example (Ship Stabilisation)

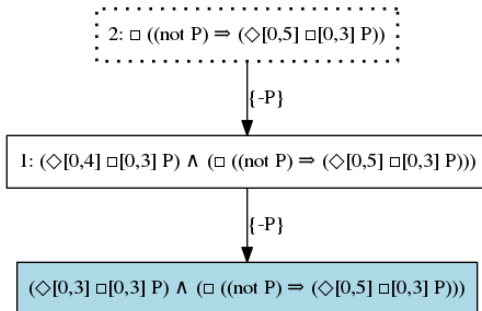
Suppose we have an autonomous ship with a landing deck. The ship attempts to stabilise itself according to the rule:

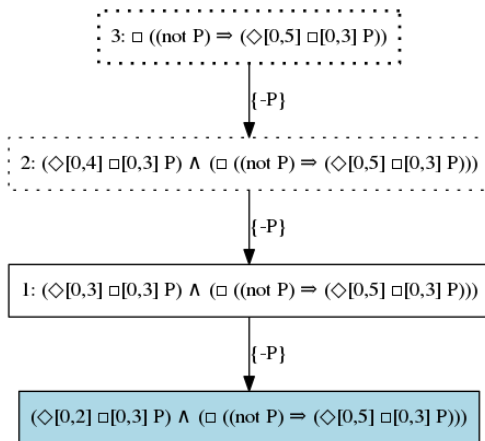
$$\Box(\neg p \rightarrow (\Diamond_{[0,5]}\Box_{[0,3]}p))$$

“Whenever the ship is unstable ($\neg p$), the ship will be stable (p) for a consecutive period of 3 minutes, within 5 minutes from having become unstable.”

$$\Box ((\text{not } P) \Rightarrow (\Diamond[0,5] \Box[0,3] P))$$



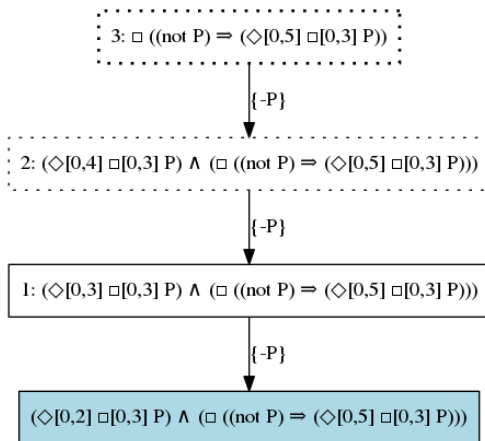


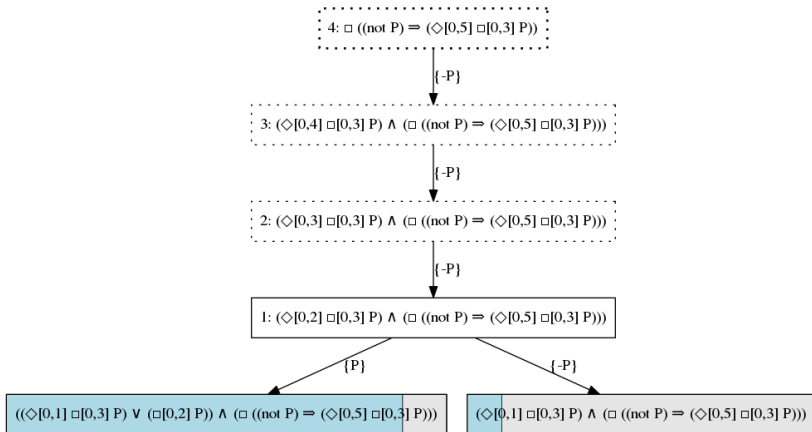


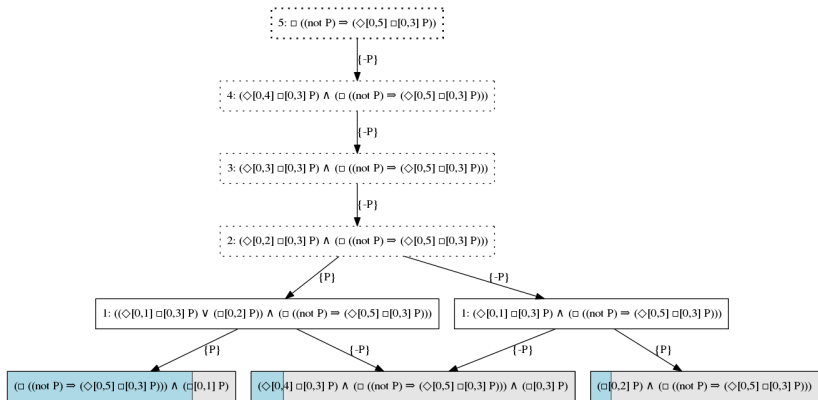
Incomplete Information

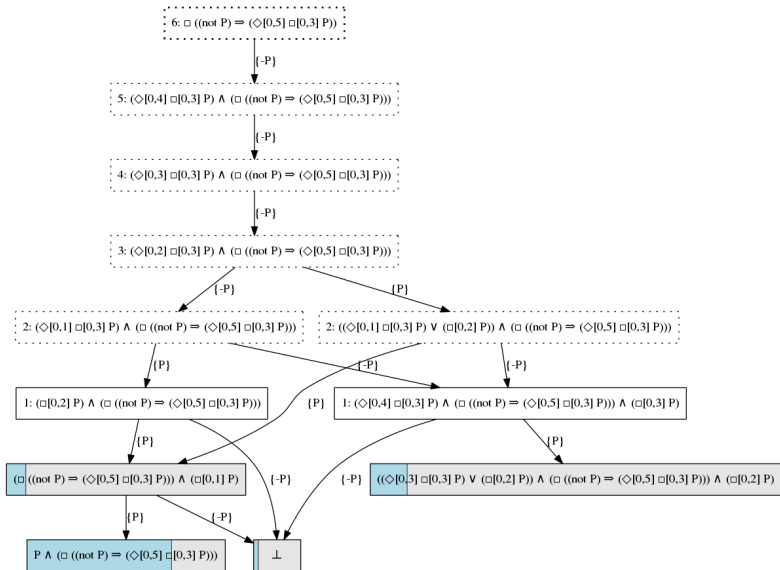
Example (Ship Stabilisation (Cont'd))

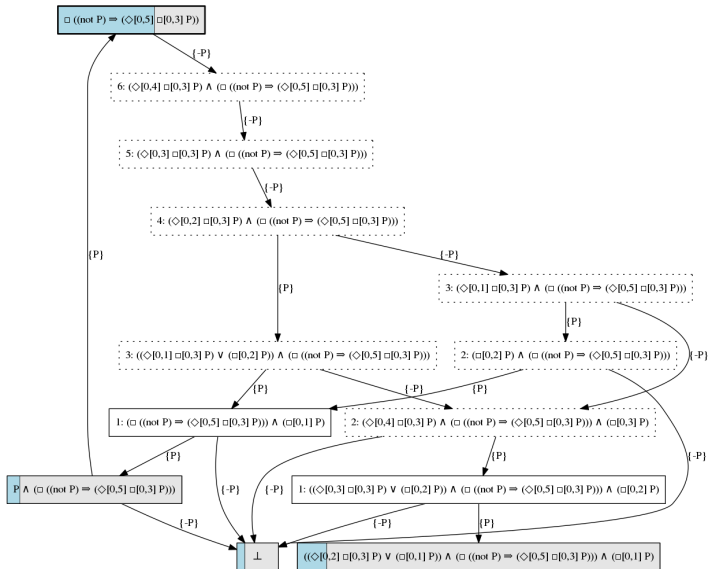
Suppose we are no longer able to measure unambiguously whether the ship is stable. Continue progression, and assume 90% stable, 10% unstable.

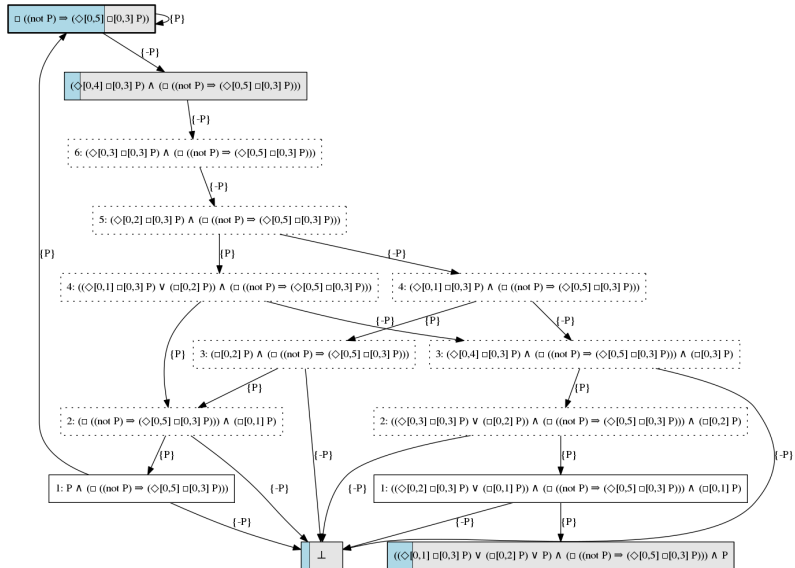


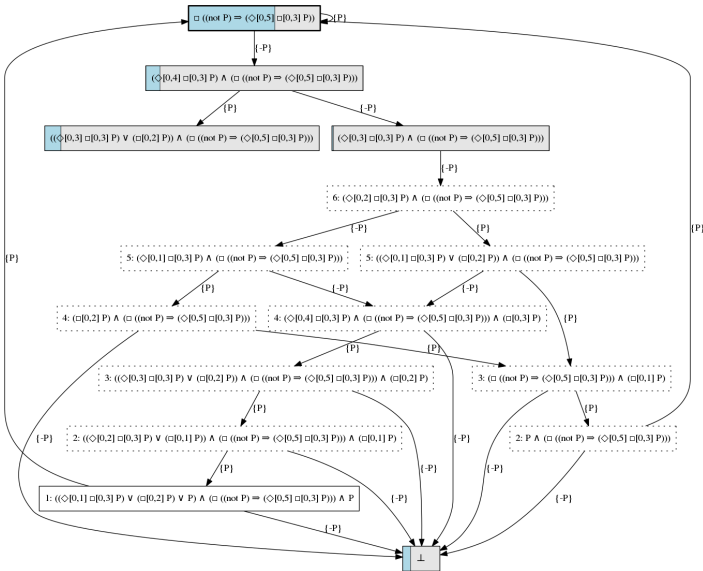


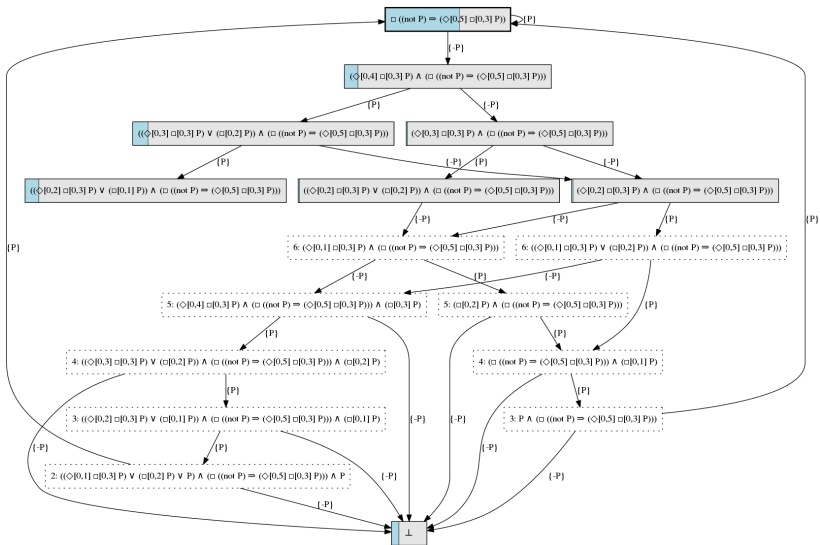












Example: Ship Stabilisation

Example (Ship Stabilisation (Cont'd))

After 10 minutes, despite incomplete sensor readings, we know:

$$Pr[\hat{\rho}, t_0 \not\models \Box(\neg p \rightarrow (\Diamond_{[0,5]}\Box_{[0,3]}p))] \geq 0.212680,$$

right now based on $m_{10}(\perp)$, *regardless* of future readings.

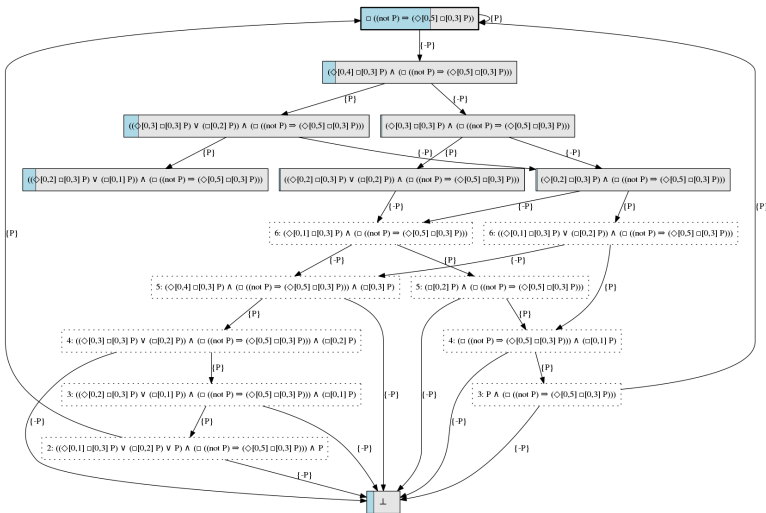
Approximate Progression

Approximate progression allows us to trade **precision** for **speed** and vice-versa:

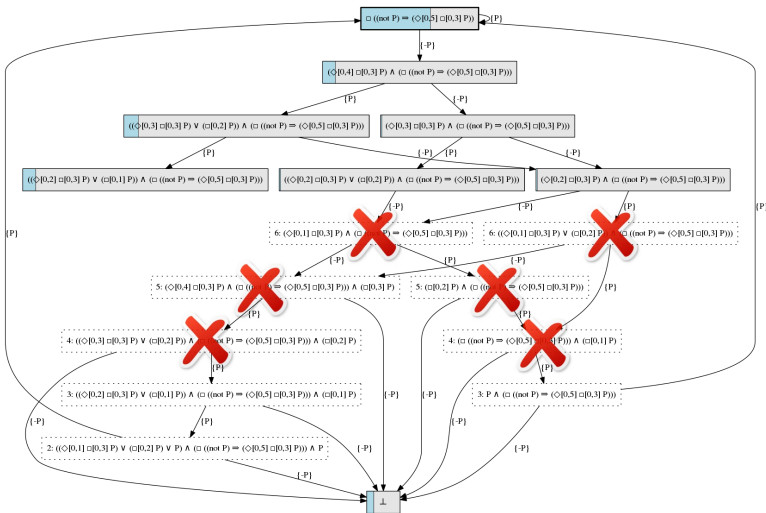
- ① Institute a MAX_AGE for formulas;
- ② Limit the size of the graph by setting a MAX_NODES bound.

We may drop nodes with probability mass, thereby **leaking** some probability mass over time.

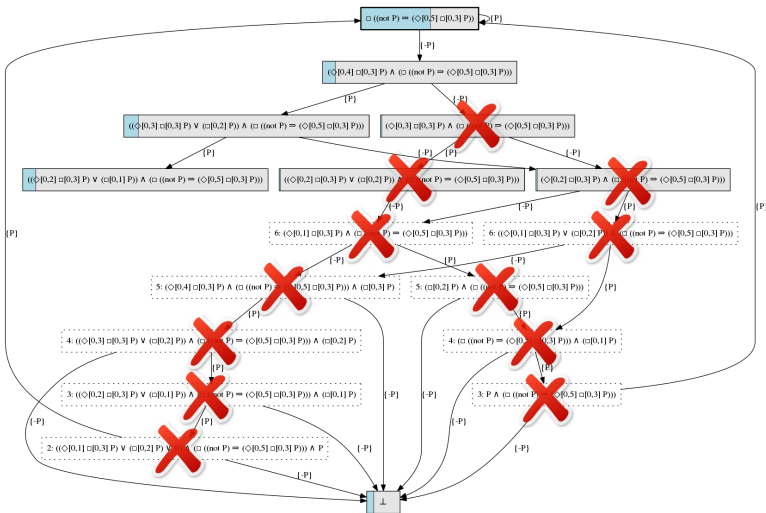
Methods to **reduce the graph size**: MAX_AGE and MAX_NODES.



Performance penalty: MAX_AGE = 3



Precision penalty: MAX_NODES = 5



Summary

Summary:

- 1 Classical progression assumes complete states;
- 2 We extended progression to handle incomplete states;
- 3 Progression graphs allow us to implicitly keep track of traces;
- 4 Approximation offers a trade-off between precision and speed.

Many interesting extensions possible!